

# Discrete Wavelet Transforms in the Large Time-Frequency Analysis Toolbox for Matlab/GNU Octave

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The discrete wavelet transform module is a recent addition to the Large Time-Frequency Analysis Toolbox (LTFAT). It provides implementations of various generalizations of Mallat's well known algorithm (iterated filterbank) such that completely general filterbank trees, dual-tree complex wavelet transforms and wavelet packets can be computed. The resulting transforms can be equivalently represented as filterbanks and analyzed as filterbank frames using fast algorithms.

CCS Concepts: • **Mathematics of computing** → **Computation of transforms; Approximation;**

General Terms: Algorithms, Documentation

Additional Key Words and Phrases: Filterbanks, finite frames, discrete wavelet transform, discrete wavelet packets, dual-tree complex wavelet transform

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## 1. INTRODUCTION

Historically, the wavelet transform and filterbanks were connected through the theory of Multiresolution Analysis (MRA) [Mallat 1989]. The MRA wavelet basis functions do not admit a closed form but they are created recursively using perfect reconstruction elementary filterbanks. MRA-based wavelet bases usually exhibit only a dyadic scale sampling of the underlying wavelet transform but the fact that they form bases in conjunction with existence of the fast Mallat's algorithm ensured a widespread adoption of a wavelet transform in such a form. For an introduction to the wavelet transform see e.g. [Strang and Nguyen 1997; Burrus et al. 2013].

Many scientists had focused on extending the MRA concept or Mallat's algorithm itself trying to overcome shortcomings of the original MRA-based wavelet transform design. Such efforts were e.g.  $M$ -band wavelet transforms [Steffen et al. 1993], wavelet packets [Coifman and Wickerhauser 1992], undecimated wavelet transforms [Holschneider et al. 1990], framelets [Daubechies et al. 2003] and the dual-tree wavelet transform [Kingsbury 2000] and others.

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After the initial enthusiasm faded, the software support started lagging behind the theory being developed and nowadays it seems that most of the once popular toolboxes are not actively maintained, or they are even no longer available for downloading. Moreover, new scientific findings in the field of discrete wavelet transforms have become scattered between many small toolboxes of varying quality, usually meant only as a proof-of-concept implementation accompanying a research paper or a book.

Several existing wavelet toolboxes are worth mentioning. In the following list, the year in brackets is the year of the last official release.

- Uvi\_Wave toolbox (1996), no official webpage [Prelicic et al. 1996], GPL
- Wavelab (2005) <http://statweb.stanford.edu/~wavelab>, custom GPL-like license
- WavBox (2012) <http://www.toolsmiths.com/Wavelet/WavBox>, commercial.
- Official Matlab Wavelet Toolbox<sup>TM</sup> (2015) <http://www.mathworks.com/products/wavelet>, commercial.

Although most of the wavelet theory leans on the theory of frames, not much attention has been paid to the frame properties of the filterbank structures that the transforms are actually computed with, like e.g. MRA filterbanks were initially meant only as a tool for computing approximate coefficients of a wavelet series. More precisely, none of the toolboxes treats any of the wavelet transforms as frames.

The discrete formalism of MRA in  $\ell^2(\mathbb{Z})$  was introduced in [Rioul 1993]; we follow these ideas (to some extent) and extend them while relying on finite frame theory in  $\mathbb{C}^L$ .

In this paper, we present a part of the Large Time-Frequency Analysis Toolbox<sup>1</sup> (LTFAT) which tries to bridge the gap in the scattered wavelet software by collecting wavelet transform generalizations, while providing a unified interface to all of them and by providing tools for analyzing their frame properties. Its aim is to provide a self-contained, modern, fast and easy to use toolbox serving as a solid base for further scientific developments.

We believe that such an effort can only be useful if the toolbox is available freely, therefore LTFAT is an open-source software licensed under a permissible GPLv3 licence and it can be obtained free of charge at <http://lftat.github.io>. GNU Octave, a free Matlab alternative, is fully supported for the same reason. In fact, LTFAT has been an official Octave-Forge package since 2013.

Apart from discrete wavelets, LTFAT covers a large collection of time-frequency transforms, an object-oriented frames framework and support for real-time audio processing directly from Matlab/GNU Octave. Moreover, computationally intensive routines are programmed duplicately in the C (backend library and MEX files) and C++ (OCT files) languages to achieve speedup. Interested readers can find more details in the overview papers [Søndergaard et al. 2012; Průša et al. 2014].

The purpose of this paper is to describe the algorithms used in the wavelets module and establish the connection to finite frames in  $\mathbb{C}^L$ . Here we list features not to be commonly found in other toolboxes:

- Free and open-source software.
- Extensible wavelet filter database.
- No restriction to critically subsampled two-channel cases.
- Support for creating completely custom filterbank trees.
- Routines for casting filterbank trees to identical non-iterated filterbanks.
- Calculation of frame bounds of any shape of wavelet filterbank trees.
- Online documentation including examples.

<sup>1</sup>The current version 2.1.1 was released on 23. 10. 2015.

- Fully supported in GNU Octave 3.8.0 and above.
- Possibility to use discrete wavelets in the frames framework which offers a common interface for most transforms in LTFAT.

To be complete, there are still areas from the wavelet theory the toolbox is lacking:

- Boundary-adjusted filters (see e.g. WaveLab).
- Advanced multidimensional transforms, as the toolbox is primarily devoted to one-dimensional signal processing.
- The discretized variant of the continuous wavelet transform with classical wavelets such as Mexican hat, Morlet wavelet etc. The reason is that the general approach to discretization does not ensure a perfect reconstruction and therefore the mathematical frame abstraction is no longer applicable.
- A “lifting” form of the basic wavelet filterbanks [Sweldens 1996]. The lifting scheme was introduced as an alternative to the basic wavelet filterbanks both as a design and a computational procedure. It might be included in future releases of LTFAT.

The rest of the paper is organized as follows. Section 1.2 establishes basic concepts of finite frame theory, sections 1.3 and 1.4 define basic properties of multirate filterbank systems, which serve for both the computation and the analysis of transform properties in the subsequent sections. Sections 2, 3, 4, 5, 6 introduce variants of wavelet type transforms currently supported in LTFAT, while section 7 wraps up with the description of an algorithm for evaluating frame properties of all the described transforms.

### 1.1. Notation

Throughout this paper, we will be working exclusively with finite-length signals denoted by a bold lower-case letter, represented as a column vector  $\mathbf{f} = [\mathbf{f}(0), \mathbf{f}(1), \dots, \mathbf{f}(L-1)]^T$  belonging to  $\mathbb{C}^L$  or  $\mathbb{R}^L$  with periodic indexing such that  $\mathbf{f}(l+kL) = \mathbf{f}(l)$  for  $l, k \in \mathbb{Z}$ . Such a treatment effectively makes the signal one period of an infinite, periodic signal, which, although it causes sometimes unnatural behaviour near the signal boundaries, makes the algorithms and the mathematical treatment clearer. Matrices will be denoted by bold capital letters, e.g.  $\mathbf{F}$ , and a conjugate transpose will be denoted as  $\mathbf{F}^*$ . The identity matrix will be denoted by  $\mathbf{I}$ . The scalar product in  $\mathbb{C}^L$  is defined as  $\langle \mathbf{x}, \mathbf{y} \rangle = \sum_{l=0}^{L-1} \mathbf{x}(l)\mathbf{y}(l)^*$  and the induced norm as  $\|\mathbf{x}\| = \sqrt{\langle \mathbf{x}, \mathbf{x} \rangle}$ . A circular convolution is  $(\mathbf{x} \circledast \mathbf{y})(n) = \sum_{l=0}^{L-1} \mathbf{x}(l)\mathbf{y}(n-l)$  for  $n \in \{0, 1, \dots, L-1\}$ . Upsampling by a factor of  $a$  is a map  $\uparrow_a: \mathbb{C}^N \rightarrow \mathbb{C}^{aN}$  such that  $(\uparrow_a \mathbf{x})(an) = \mathbf{x}(n)$  for  $n \in \{0, 1, \dots, N-1\}$  and zero at the other positions. Downsampling by a factor  $a$  is a map  $\downarrow_a: \mathbb{C}^{aN} \rightarrow \mathbb{C}^N$  such that  $(\downarrow_a \mathbf{x})(n) = \mathbf{x}(an)$  for  $n \in \{0, 1, \dots, N-1\}$ . Vector reflection and conjugation (involution) will be denoted by  $\bar{\mathbf{x}}(l) = \mathbf{x}(-l)^*$  for  $l \in \{0, \dots, L-1\}$  assuming periodic indexing. The discrete Fourier transform of  $\mathbf{f}$  will be defined as  $\hat{\mathbf{f}}(n) = (\mathcal{F}\mathbf{f})(n) = \sum_{l=0}^{L-1} \mathbf{f}(l) \exp(-i2\pi nl/L)$  and the inverse transform as  $\mathbf{f}(l) = (\mathcal{F}\hat{\mathbf{f}})(l) = 1/L \sum_{n=0}^{L-1} \hat{\mathbf{f}}(n) \exp(i2\pi nl/L)$ . The positive remainder of division  $a/b$  is denoted as  $(a \bmod b)$  and, finally,  $\text{lcm}(a_0, a_1, \dots, a_{M-1})$  denotes the least common multiple of  $a_0, \dots, a_{M-1}$  such that  $\text{lcm}(a_0, a_1, \dots, a_{M-1}) = \text{lcm}(\text{lcm}(a_0, a_1), a_2), \dots$

*Toolbox conventions.* In the actual implementation in Matlab/Octave, the signal does not have to be a column vector. When a single row vector is used, it is transposed internally, for matrices the operation is *broadcast*, i.e. the operation is applied to each column or optionally to each row, if the function supports specification of the *dim* parameter.

The LTFAT toolbox functions will be referred to in a typewriter font e.g. `funname`. The complete description of functions is available at the documentation webpage <http://lftfat.github.io/doc/> or via running `help funname`.

A custom function input argument parser `lftfatarghelper`<sup>2</sup> was included in LTFAT to make the function calls more versatile and clear. It supports positional arguments, key-value pairs and string flag groups and it allows specifying default values for them. Such features are not contained in Matlab/Octave by default.

*Algorithms conventions.* In the description of the algorithms we make use of dynamic array constructs such as a First-In First-Out (FIFO) queue and Last-In Last-Out (LIFO) stack. Operations defined on the queue will be *enqueue* meaning appending an element as the last one in the queue and *dequeue* meaning removal of the first element from the queue. Operations with a stack are *push* and *pop* performing addition and removal of the last element respectively.

## 1.2. Finite frames

The frame theory encompasses all linear, perfect reconstruction transforms. In the finite setting [Balazs 2008; Kovačević and Chebira 2008; Casazza and Kutyniok 2013], a frame for  $\mathbb{C}^L$  is a set of  $\Lambda$  vectors in  $\mathbb{C}^L$  which can be represented as columns of an  $L \times \Lambda$ , ( $L \leq \Lambda$ ) full row rank matrix  $\mathbf{F}$ . We further restrict  $\Lambda < \infty$  to exclude infinite matrices. Every frame  $\mathbf{F}$  is equipped with frame bounds  $0 < A \leq B < \infty \in \mathbb{R}$  with

$$A \|\mathbf{f}\|^2 \leq \|\mathbf{F}^* \mathbf{f}\|^2 \leq B \|\mathbf{f}\|^2 \quad (1)$$

for all  $\mathbf{f} \in \mathbb{C}^L$ . The optimal bounds are the tightest possible in (1) and correspond to squares of the min and max singular values of  $\mathbf{F}$  respectively. In the following text  $A, B$  will denote exclusively the optimal frame bounds. A frame is called *overcomplete* if  $\Lambda > L$ .

The fundamental operators associated with frames are the *analysis* operator  $\mathbf{F}^* : \mathbb{C}^L \rightarrow \mathbb{C}^\Lambda$ , and the *synthesis* operator  $\mathbf{F} : \mathbb{C}^\Lambda \rightarrow \mathbb{C}^L$ . In the finite setting, they can both be represented as matrices being conjugate transpose of each other. An invertible *frame operator* ( $L \times L$  matrix) is formed as concatenation of the analysis and the synthesis operators,  $\mathbf{F}\mathbf{F}^*$ . The frame bounds  $A, B$  are also equal to min and max eigenvalues of the frame operator respectively. Each frame is associated with a unique *canonical dual* frame such that  $\mathbf{F}_d = (\mathbf{F}\mathbf{F}^*)^{-1} \mathbf{F}$  and  $\mathbf{F}_d \mathbf{F}^* = \mathbf{F}\mathbf{F}_d^* = \mathbf{I}$  ( $\mathbf{F}_d^*$  coincides with the right inverse of  $\mathbf{F}$ ).

If a frame is not overcomplete, it is an orthonormal basis if  $\mathbf{F}\mathbf{F}^* = \mathbf{I}$  and a general (Riesz) basis otherwise. In such cases, the signal representation in the frame coordinates  $\mathbf{f} = \mathbf{F}\mathbf{c}$  is unique with  $\mathbf{c} = \mathbf{F}^{-1} \mathbf{f}$  which simplifies to  $\mathbf{c} = \mathbf{F}^* \mathbf{f}$  in the orthonormal basis case. The canonical dual frame is the only one possible dual frame. However, when a frame is overcomplete, the solution to  $\mathbf{f} = \mathbf{F}\mathbf{c}$  is not unique. All possible solutions can be parametrized as  $\mathbf{c} = \mathbf{F}_d^* \mathbf{f} + (\mathbf{I} - \mathbf{F}_d^* \mathbf{F}) \mathbf{w}$  for an arbitrary vector  $\mathbf{w} \in \mathbb{C}^\Lambda$ ; the  $(\mathbf{I} - \mathbf{F}_d^* \mathbf{F})$  part denotes projection onto the null space of  $\mathbf{F}$ . The minimum energy solution is simply  $\mathbf{c} = \mathbf{F}_d^* \mathbf{f}$ .

A frame is called *tight* if  $A = B$  and  $\mathbf{F}\mathbf{F}^* = A\mathbf{I}$  and *Parseval tight* if in addition  $A = 1$ . Any tight frame can be normalized such that it becomes Parseval tight. If the frame is tight, the calculation of the canonical dual frame reduces to  $\mathbf{F}_d = \frac{1}{A} \mathbf{F}$ . A frame admits a *painless* expansion [Daubechies et al. 1986] if the frame operator is diagonal and therefore easily invertible. When a frame is not tight, the frame operator is seldom diagonal in the original domain but it can become diagonal, or at least structured conveniently in a different domain. Applying a unitary operator  $\mathbf{U}$  to each frame element

<sup>2</sup>Since version 2.1.0 `lftfatarghelper` is also available as a MEX function.

results in an *unitary isomorphic* frame [Casazza and Kutyniok 2013]  $\mathbf{F}_{\text{ui}} = \mathbf{U}\mathbf{F}$  for which  $\mathbf{F}_{\text{ui}}^* \mathbf{F}_{\text{ui}} = \mathbf{F}^* \mathbf{F}$  and therefore the frame bounds of  $\mathbf{F}$  and  $\mathbf{F}_{\text{ui}}$  are equal. In the latter sections we will use the (unitary) discrete Fourier transform as  $\mathbf{U}$ .

In general, there is an infinite number of other dual frames  $\mathbf{F}_{\text{dd}}$  such that  $\mathbf{F}_{\text{dd}} \mathbf{F}^* = \mathbf{F} \mathbf{F}_{\text{dd}}^* = \mathbf{I}$ . This is an important property of overcomplete frames as the canonical dual frame might not keep the same structure (and therefore admit a fast algorithm) in contrary to some other dual frame. On the other hand, the canonical dual frame plays an important role when the coefficients are modified and one needs to make a projection onto the range space of  $\mathbf{F}^*$  or null space of  $\mathbf{F}$ . Such operations are required in some iterative schemes e.g. in the Griffin-Lim algorithm for a phase retrieval [Griffin and Lim 1984] or in the basis pursuit using the Alternating Direction Method of Multipliers [Boyd et al. 2011]. The canonical dual frames are also important in the context of frame multipliers [Balazs 2007].

When the canonical dual frame does not admit a fast algorithm and constructing it explicitly would not be feasible, iterative algorithms exist which calculate the inverse of a frame operator. As an example we mention Neumann series expansion and others (see e.g. [Gröchenig 1993]) which can work without explicit matrices with or without the knowledge of the frame bounds.

*Toolbox Conventions.* In LTFAT, no matrix is explicitly created as it would not be feasible for large  $L$  (tens of thousands and more). Instead, fast algorithms are exploited to perform the actual matrix-vector multiplications and inversions. All frames that will be considered admit a fast algorithm to be used in place of their operators.

### 1.3. Filterbanks

Filters  $\mathbf{g} \in \mathbb{C}^L$  considered in this paper will be finite impulse response (FIR), which means that the actual filter coefficients  $\mathbf{g}_{\text{supp}} \in \mathbb{C}^{L_{\mathbf{g}}}$  are padded with zeros to  $L$  (assuming  $L_{\mathbf{g}} \leq L$ ) and subsequently circularly shifted upwards by  $d$  samples, where  $d$  is a filter-specific initial “offset”. We will use a filter generating function  $\mathbf{g} = g(L)$  with  $(g(L))((n-d) \bmod L) = \mathbf{g}_{\text{supp}}(n)$  for  $n \in \{0, 1, \dots, L_{\mathbf{g}} - 1\}$  and producing zeros at other indices. We will refer to filters without  $L$  yet specified as  $g$ . When such a filter is used in a convolution, we assume  $L$  to be deduced from the other argument. A generating function for the upsampled version of such a filter will be defined as  $((\uparrow_a g)(L))((a(n-d)) \bmod L) = \mathbf{g}_{\text{supp}}(n)$  for  $n$  as before and the following holds:  $(\uparrow_a g)(L) = \uparrow_a g(L/a)$ . A filterbank is a collection of impulse responses  $g_m$  for  $m \in \{0, 1, \dots, M-1\}$ , followed by channel dependent integer subsampling factors  $a_m$ . A filterbank is called *uniform* if  $a_m = \dots = a_0 = a$ . The filterbank produces length  $N_m$  *subband* coefficients  $\mathbf{c}_m \in \mathbb{C}^{N_m}$  such that  $L = k \text{lcm}(a_0, a_1, \dots, a_{M-1})$  for some integer  $k \in \mathbb{N}$ :

$$\mathbf{c}_m(n) = (\downarrow_{a_m}(\mathbf{f} \circledast g_m))(n) = \sum_{l=0}^{L-1} \mathbf{f}(l) \mathbf{g}_m(a_m n - l), \quad (2)$$

for  $m \in \{0, 1, \dots, M-1\}$ ,  $N_m = L/a_m$  and  $n \in \{0, 1, \dots, N_m - 1\}$ . Formula (2) corresponds to an analysis operator of a filterbank frame defined by  $\mathbf{g}_m, a_m$  as

$$\boldsymbol{\psi}_{m,n}(l) = \{\overline{\mathbf{g}_m}(l - a_m n)\}, \quad (3)$$

for  $l \in \{0, 1, \dots, L-1\}$  and  $n$  and  $m$  as above.

The adjoint operation (the synthesis operator) is given by

$$\tilde{\mathbf{f}}(l) = \sum_{m=0}^{M-1} ((\uparrow_{a_m} \mathbf{c}_m) \circledast \overline{g_m})(l) = \sum_{m=0}^{M-1} \sum_{n=0}^{N_m-1} \mathbf{c}_m(n) \overline{\mathbf{g}_m}(a_m n - l), \quad (4)$$

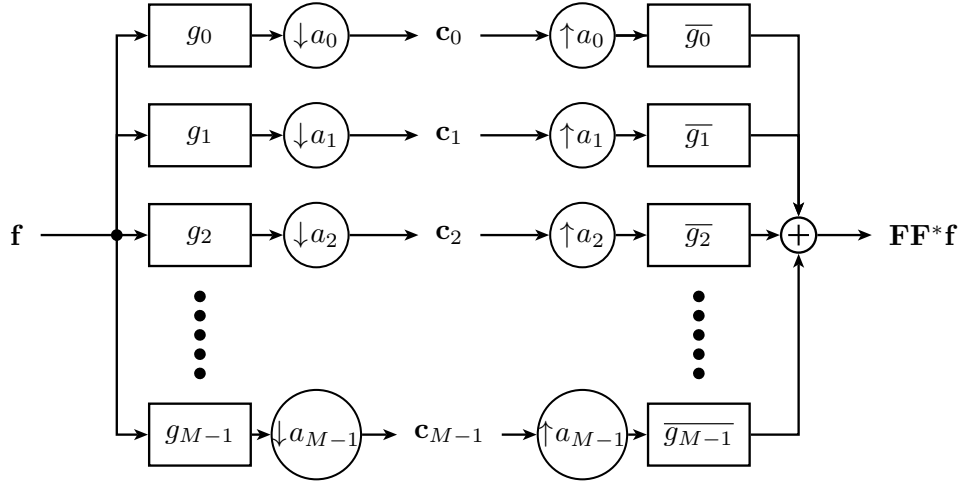


Fig. 1. Analysis and synthesis operators of a filterbank frame given as a set of filters  $g_m$  and subsampling factors  $a_m$ .

for  $l$  and  $N_m$  as above. In order to get a perfect reconstruction, one has to combine the analysis operator and the synthesis operator of dual filterbank frames.

Uniform filterbanks were studied in the context of frames in [Cvetkovic and Vetterli 1998; Bölcskei et al. 2002], the results were generalized to the  $L$ -periodic sequences in [Fickus et al. 2013], which fits in the setting assumed in this paper.

*Toolbox conventions.* Since the set of admissible  $L$  is restricted to  $L = k \text{ lcm}(a_0, a_1, \dots, a_{M-1})$ ,  $k \in \mathbb{N}$  the input signal of length  $L_s$  is padded with zeros to the proper length by default.

From a signal processing point of view, equations (2), (4) can be evaluated quickly using a time-domain polyphase implementation when working with short FIR filters or by a fast convolution using FFTs (see e.g. [Vetterli et al. 2014]). Both algorithms are well known and implemented in the computational functions called from `filterbank` and `ifilterbank`. The FIR filters are passed as a cell array of structures with fields `.h` and `.offset` representing  $g_{\text{supp}}$  and  $-d$  mentioned earlier respectively.

#### 1.4. Multirate Identity Rules

Multirate identity rules [Vaidyanathan 1993] refer to transformations of a collection of filters and rate converters which does not change the overall transfer function.

In particular, a chain of  $J$  filters and downsamplers  $g_k, a_k$  for  $k \in \{0, 1, \dots, J-1\}$  can be transformed to a single “identical” filter  $g_{\text{tot}}$  followed by a downsampling factor  $a_{\text{tot}}$  using Algorithm 1 (see Fig. 2).

## 2. THE DISCRETE WAVELET TRANSFORM

The Discrete Wavelet Transform (DWT) based on MRA, considered in this paper, is most often associated with Mallat’s algorithm [Mallat 1989] and it is then termed the Fast Wavelet Transform (FWT). It consists of an iterative application of a critically subsampled basic two-channel wavelet filterbank. The algorithm was described and analyzed by many authors (see e.g. [Shensa 1992]). The most appealing property of the algorithm is its linear complexity. The main property of the transform is a one filter per octave frequency resolution and a logarithmic overall frequency scale coming from the dyadic sampling of translation and scale of the underlying continuous transform.

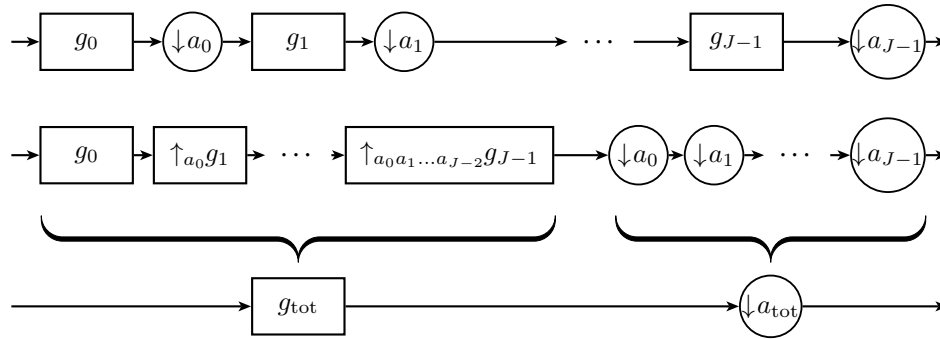


Fig. 2. Chain of filters  $g_k$  and downsampling factors  $a_k$  (top row) and the identical filter  $g_{tot}$  and downsampling factor  $a_{tot}$  (bottom row). The middle row shows an intermediate state after moving all subsampling factors to the rightmost position.

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**ALGORITHM 1:** Create a filter identical to a chain of filters and downsamplers

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**Input:** Chain of filters and downsamplers  $g_k, a_k$  for  $k \in \{0, 1, \dots, J-1\}$ .

**Output:** Identical filter  $g_{tot}$  and downsampling factor  $a_{tot}$ .

Define:  $a_{part}(k) = \begin{cases} \prod_{n=0}^k a_n & \text{if } k \geq 0 \\ 1 & \text{otherwise} \end{cases}$  ;

Define:  $L_g(k) = L_{g_0} + \sum_{n=1}^k (L_{g_n} - 1) a_{part}(n-1)$ ;

Define:  $d(k) = d_0 + \sum_{n=1}^k d_n a_{part}(n-1)$ ;

$g_{tot} = g_0$ ;

**for**  $k \leftarrow 1$  **to**  $J-1$  **do**

$g_{tot} \leftarrow (\uparrow_{a_{part}(k-1)} g_k) \otimes g_{tot}(L_g(k)); g_{tot} \leftarrow$  from  $g_{tot}$  and  $d(k)$ ;

**end**

$d_{tot} \leftarrow d(J-1); a_{tot} \leftarrow a_{part}(J-1)$ ;

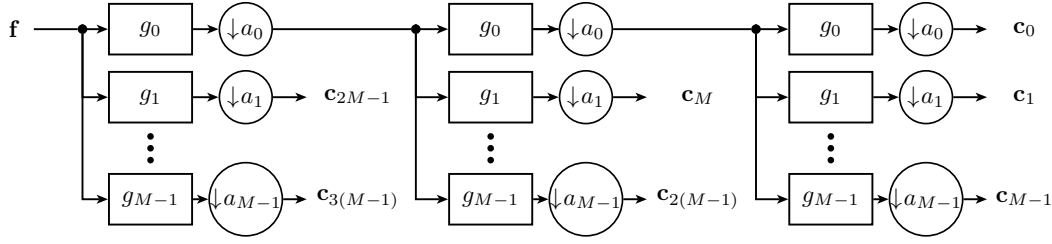
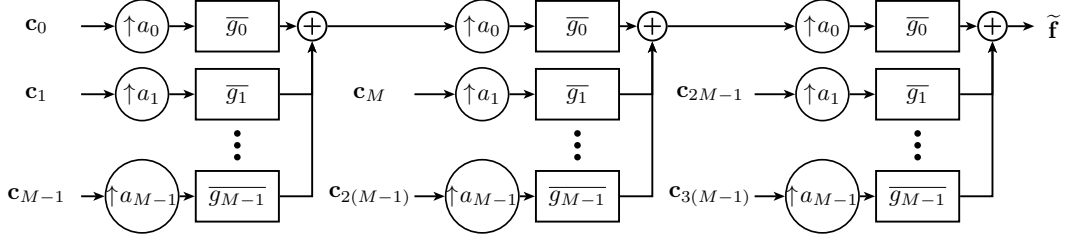
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The  $M$ -band wavelet bases [Steffen et al. 1993] generalize DWT such that they exhibit  $M-1$  linear bands per octave as  $M$  filters are used in each of the iterations. Further generalizations came from the idea of using an overcomplete filterbank frame in place of the basic wavelet filterbank (see references below, in the list of supported wavelet filters).

All the classes mentioned share the same overall iterated filterbank structure and therefore they are treated as a single type of transform. In general, a basic wavelet filterbank consists of  $M$  filters  $g_m, a_m$  with  $g_0$  denoting the low-pass filter.

The fast algorithms for calculating the output of analysis and synthesis operators are given by Alg. 2 and Alg. 3 respectively. Alg. 2 for  $J=3$  iterations is depicted in Fig. 3. In order to get a perfect reconstruction, dual basic wavelet filters must be used in one of the operators.

The algorithm effectively follows a filterbank tree, which can be equivalently represented by a non-uniform filterbank using Alg. 1 such that each output is associated with a chain of filters and downsamplers in a path from the root node to the output. Although the identical filterbank admits a less effective algorithm in general, it allows investigating properties of the resulting subbands  $c_k$  like frequency bands and group delay and it also determines the effective frame vectors as in (3). In the usual cases when the frequency responses of basic filterbank filters cover distinct frequency bands ordered by increasing center frequency, the identical filterbank frequency response is ordered in the same way.

Fig. 3. Iterated filterbank structure with  $J = 3$ .Fig. 4. Adjoint filterbank structure with  $J = 3$ .**ALGORITHM 2:** Fast wavelet analysis

**Input:** Input signal  $f \in \mathbb{C}^{L_s}$ , basic wavelet filters  $g_m, a_m$  for  $m \in \{0, 1, \dots, M-1\}$ , number of levels  $J$ .

**Output:** Coefficients  $c_k$  for  $k \in \{0, 1, \dots, J(M-1)\}$ .

```

1  $a \leftarrow f$  and pad with zeros to  $L$ , the next integer multiple of  $a_0^J$ ;
2  $k \leftarrow J(M-1)$ ;
3 for  $j \leftarrow 0$  to  $J-1$  do
4   for  $m \leftarrow M-1$  to  $1$  do
5      $c_k \leftarrow \downarrow_{a_m}(a \otimes g_m)$ ;
6      $k \leftarrow k-1$ ;
7   end
8    $a \leftarrow \downarrow_{a_0}(a \otimes g_0)$ ;
9 end
10  $c_0 \leftarrow a$ ;

```

*Toolbox conventions.* The routines implementing Algs. 2 and 3 are called `fwt` and `ifwt` respectively. Wavelet filters are stored in a direct filterbank form as defined in sec. 1.3. In fact, a pair of (dual) basic wavelet filterbanks is always stored (or computed) and the filterbanks differ only in a non-tight frame case. The basic wavelet filter impulse responses are stored in, or generated by functions with `wfilt_` prefix. LTFAT currently includes the classical ones i.e. Daubechies filters `db`, symlets `sym`, coiflets `coif`, biorthogonal spline wavelet filters `spline` [Daubechies 1992] and some less common filters e.g. near-orthogonal symmetric and symmetric near-orthogonal `symorth` [Abdelnour and Selesnick 2004]. Additional supported wavelet filters are:  $M$ -band wavelets `cmband`, `mband`<sup>3</sup>, `algband` [Gopinath and Burrus 1995; Alkin and Caglar 1995; Lin et al.

<sup>3</sup>This is technically not a proper basic wavelet filterbank as the low-pass filter does not fulfill the regularity condition, but nevertheless, for a finite number of transform levels, it creates an orthonormal basis.



**ALGORITHM 3:** Fast wavelet synthesis

**Input:** Coefficients  $c_k$  for  $k \in \{0, 1, \dots, J(M-1)\}$ , basic wavelet filters  $g_m, a_m$  for  $m \in \{0, 1, \dots, M-1\}$ , number of levels  $J$ , original input signal length  $L_s$ .

**Output:** Reconstructed signal  $\tilde{\mathbf{f}} \in \mathbb{C}^{L_s}$ .

```

1  $\mathbf{a} \leftarrow \mathbf{c}_0$ ;
2  $k \leftarrow 1$ ;
3 for  $j \leftarrow J-1$  to 0 do
4    $\mathbf{a} \leftarrow (\uparrow_{a_0} \mathbf{a}) \otimes \overline{g_0}$ ;
5   for  $m \leftarrow 1$  to  $M-1$  do
6      $\mathbf{a} \leftarrow \mathbf{a} + (\uparrow_{a_m} \mathbf{c}_k) \otimes \overline{g_m}$ ;
7      $k \leftarrow k+1$ ;
8   end
9 end
10  $\tilde{\mathbf{f}} \leftarrow \mathbf{a}(0, 1, \dots, L_s-1)$ ;

```

2006], wavelet filters for constructing overcomplete tight or general filterbank frames `dden`, `symdden`, `symtight`, `hden`, `dgrid`, `symds`, [Selesnick 2001; Selesnick and Abdelnour 2004; Abdelnour and Selesnick 2005; Selesnick 2006; Abdelnour 2007; 2012].

The initial position ( $d$  from sec. 1.3) of the basic filters is usually not addressed in the literature since only the relative alignment of the filters makes a difference from the point of view of the MRA theory. However it might become a source of confusion in the finite setting with periodic boundary conditions as the frame vectors associated with different levels of DWT can be supported at distinct time positions causing misalignment of the coefficient subbands. Most of the basic wavelet filters are not symmetric and therefore do not admit a distinctive “center” point. Some toolboxes simply do not treat this detail, others employ ad-hoc methods for finding the center point while keeping the relative filter alignment intact but the methods might be suboptimal in some cases. In LTFAT, the default position of the basic filters was hand-tuned to minimize such ambiguity.

### 3. GENERAL FILTERBANK TREES AND WAVELET PACKETS

By a filterbank tree we understand a tree-shaped, oriented (no backwards connections) interconnection of basic wavelet filterbanks. DWT can be regarded as a special shape of a filterbank tree where only the low-pass output is decomposed up to depth  $J$  and all the nodes contain the same basic filterbank. In general filterbank trees, any further branching is possible and even the filterbanks in the nodes might differ; see Fig. 5 for a filterbank tree example. This results in a modular division of the frequency band. Alg. 4 acts as the analysis operator of a filterbank tree frame, whereas Alg. 5 acts as the synthesis operator. In order to obtain a perfect reconstruction, the corresponding nodes in the analysis tree and the synthesis tree have to be dual filterbanks.

Wavelet packets [Coifman and Wickerhauser 1992] share the same tree shape with the general filterbank trees, but in contrary to the common wavelet transform, all the intermediate outputs of all nodes are kept as additional subbands. In the literature, there is usually no distinction between wavelet filterbank trees and wavelet packets, but we have made the distinction to emphasize the slightly different algorithms and to be able to investigate different frame properties of both representations.

Since the highly redundant wavelet packet representation uses the same filters as the wavelet filterbank trees, it is necessary to introduce scaling of all intermediate subbands (or filterbanks which produce them) in the analysis and/or synthesis operation, in order to achieve perfect reconstruction. Three types of scaling are supported: no scaling, by  $1/2$  or by  $1/\sqrt{2}$ . Wavelet packet analysis is performed according to Alg. 6

and the synthesis according to Alg. 7. Again, in order to obtain a perfect reconstruction, dual filterbanks have to be used in the corresponding nodes during analysis and synthesis in addition to the compatible scaling of the intermediate subbands.

---

**ALGORITHM 4:** Filterbank tree analysis
 

---

**Input:** Input signal  $\mathbf{f} \in \mathbb{C}^{L_s}$ , filterbank tree  $\mathcal{T}$  with  $N$  nodes in the breadth-first order where the  $n$ -th node consists of filterbank  $g_m^n, a_m^n$  for  $m \in \{0, 1, \dots, M_n - 1\}$ . Total number of unconnected outputs of all nodes is  $K$ .

**Output:** Coefficients  $\mathbf{c}_k$  for  $k \in \{0, 1, \dots, K - 1\}$  in the natural ordering.

Determine the total subsampling factors for all outputs  $a_{\text{tot},k}$  according to Alg. 1;

$\mathbf{a} \leftarrow \mathbf{f}$  and pad with zeros to  $L$ , the next integer multiple of  $\text{lcm}(a_{\text{tot},0}, a_{\text{tot},1}, \dots, a_{\text{tot},K-1})$ ;

$kIdMap \leftarrow$  Alg. 20 with  $\mathcal{T}$ ;

$k \leftarrow 0$ ;

$aFifo \leftarrow \{\}$ ; % an empty FIFO queue of vectors

enqueue  $\mathbf{a}$  to  $aFifo$ ;

**for**  $n \leftarrow 0$  **to**  $N - 1$  **do**

$\mathbf{a} \leftarrow$  dequeue from  $aFifo$ ;

**for**  $m \leftarrow 0$  **to**  $M_n - 1$  **do**

$\mathbf{b} \leftarrow \downarrow_{a_m^n}(\mathbf{a} \otimes g_m^n)$ ;

**if** node  $n$  has a next node connected at index  $m$  **then**

            enqueue  $\mathbf{b}$  to  $aFifo$ ;

**else**

$\mathbf{c}_{kIdMap(k)} \leftarrow \mathbf{b}$ ;

$k \leftarrow k + 1$ ;

**end**

**end**

**end**

---

As in the case of DWT in sec. 2, the filterbank tree can be transformed to an identical filterbank using Alg. 1 for each of the paths from the root to the  $k$ -th output of the tree. The resulting frequency bands are in the *natural* order (see sec. 3.1).

The tree can be built in a completely custom way by using functions `wfbtinit`, `wfbtput` and `wfbtremove` (see examples in the documentation of the respective functions). There are two predefined tree shapes: DWT tree ('dwt' flag) and full tree ('full' flag) decomposition up to a specified depth  $J$  repeating the same basic filterbank in each node. Moreover, using a trick due to [Bayram and Selesnick 2008] allows creating a  $2^k$ -band wavelet transform from a basic two channel elementary filterbank by using  $k$  level full decomposition trees as levels of a  $2^k$ -band DWT filterbank tree. The symbol  $J$  then denotes the number of levels of the  $2^k$ -band transform. Such a filterbank tree can be created by using 'doubleband' for  $k = 2$ , 'quadband' ( $k = 3$ ) and 'octaband' ( $k = 4$ ) flags in the computational routines.

A traditional way of creating a particular filterbank tree adapted to an input signal is based on evaluating a cost function of wavelet packet subbands. The procedure is referred to the best-basis selection [Wickerhauser 1994], which was generalized to non-additive cost functions in [Taswell 1994]. Both cases are addressed in Alg. 8. It is restricted to orthonormal basic filterbanks. Any filterbank subtree chosen by the algorithm forms an orthonormal basis.

**ALGORITHM 5:** Filterbank tree synthesis

**Input:** Coefficients  $c_k$  for  $k \in \{0, 1, \dots, K-1\}$ , filterbank tree  $\mathcal{T}$  with  $N$  nodes in the breadth-first order where the  $n$ -th node consists of filterbank  $g_m^n, a_m^n$  for  $m \in \{0, 1, \dots, M_n-1\}$  and the number of all unconnected outputs of all nodes is  $K$ , original input signal length  $L_s$ .

**Output:** Reconstructed signal  $\tilde{\mathbf{f}} \in \mathbb{C}^{L_s}$ .

```

kIdMap ← Alg. 20 with  $\mathcal{T}$ ;
k ← K - 1;
aFifo ← {}; % an empty FIFO queue of vectors
for n ← N - 1 to 0 do
  b ← 0;
  for m ← M_n - 1 to 0 do
    if node n has a next node connected at index m then
      a ← dequeue from aFifo;
    else
      a ← ckIdMap(k);
      k ← k - 1;
    end
    b ← b +  $\uparrow_{a_m^n}(\mathbf{a} \otimes \overline{g_m^n})$ ;
  end
  enqueue b to aFifo;
end
 $\tilde{\mathbf{f}} \leftarrow \mathbf{b}(0, 1, \dots, L_s - 1)$ ;

```

**ALGORITHM 6:** Wavelet packet analysis

**Input:** Input signal  $\mathbf{f} \in \mathbb{C}^{L_s}$ , filterbank tree  $\mathcal{T}$  with  $N$  nodes in the breadth-first order, where the  $n$ -th node consist of a filterbank  $g_m^n, a_m^n$  for  $m \in \{0, 1, \dots, M_n-1\}$ , scaling of intermediate subbands  $s \in \{1, 1/\sqrt{2}, 1/2\}$ .

**Output:** Coefficients  $c_k$  for  $k \in \{0, 1, \dots, -1 + \sum M_n\}$ , stack of indexes of nodes *inIdStack*.

Determine the total subsampling factors for all outputs  $a_{\text{tot},k}$  according to Alg. 1;

$\mathbf{a} \leftarrow \mathbf{f}$  and pad with zeros to  $L$ , the next integer multiple of  $\text{lcm}(a_{\text{tot},0}, a_{\text{tot},1}, \dots, a_{\text{tot},K-1})$ ;

$k = 0$ ;

*inIdFifo* = {}; % an empty FIFO queue of indices

*inIdStack* = {}; % an empty LIFO stack of indices

```

for n ← 0 to N - 1 do
  for m ← 0 to M_n - 1 do
    ck ←  $\downarrow_{a_m^n}(\mathbf{a} \otimes g_m^n)$ ;
    if node n has a next node connected at m then
      ck ← s ck;
      enqueue k to inIdFifo;
      push k to inIdStack;
    end
    k = k + 1;
  end
  if inIdFifo is not empty then
    knext ← dequeue from inIdFifo;
    a = cknext;
  end
end

```

**ALGORITHM 7: Wavelet packet synthesis**

**Input:** Coefficients  $c_k$  for  $k \in \{0, 1, \dots, -1 + \sum M_n\}$ , filterbank tree  $\mathcal{T}$  with  $N$  nodes in a breadth-first order, where  $n$ -th node consists of filterbank  $g_m^n, a_m^n$  for  $m \in \{0, 1, \dots, M_n - 1\}$ , stack of indexes  $inIdStack$ , scaling of intermediate subbands  $s \in \{1/2, 1/\sqrt{2}, 1\}$ , original input length  $L_s$ .

**Output:** Reconstructed signal  $\tilde{f} \in \mathbb{C}^{L_s}$ .

```

 $k = -1 + \sum M_n$ ;
for  $n \leftarrow N - 1$  to 0 do
  if  $n \neq 0$  then % not a root node
     $id$  pop from  $inIdStack$ ;
     $a \leftarrow c_{id}$ ;
  else
     $a \leftarrow 0$ ;
  end
  for  $m \leftarrow M_n - 1$  to 0 do
     $a \leftarrow a + (\uparrow_{a_m^n} c_k) \otimes g_m^n$ ;
     $k \leftarrow k - 1$ ;
  end
  if  $n \neq 0$  then  $c_{id} \leftarrow sa$ ;
  ;
end
 $\tilde{f} \leftarrow a(0, 1, \dots, L_s - 1)$ ;

```

*Toolbox conventions.* The filterbank trees are stored in an abstract data structure behaving like a tree. It consists of an array of nodes containing basic filterbanks from sec. 2 and two arrays defining indexes of parent and children nodes<sup>4</sup>.

The functions implementing Alg. 4 and Alg. 5 are `wfbt` and `iwfbt` respectively. Alg. 6 and Alg. 7 are implemented in `wpfbt` and `iwfbt` respectively. Function `wpbst` implements the best basis selection (algorithm 8).

### 3.1. Natural (Paley) and Frequency ordering of subbands

A filterbank tree's subbands, as obtained from Alg. 4 or 6, are not ordered according to the frequency. For a full decomposition tree and nodes containing two half-band filters, this ordering is called *natural* or *Paley*. In such a case, the subbands can be shuffled to a proper frequency ordering by rewriting the subband indexes in Gray binary code. A general way of handling the frequency ordering of subbands valid for any decomposition tree shape involves reversing the order of filters in some nodes as it is described in Alg. 9. Applying the algorithm for the second time changes the frequency subband ordering back to natural. The algorithm comes from [Wickerhauser 1994], originally designed to work with filterbanks with two filters only and it was modified to correctly treat basic wavelet filterbanks with odd numbers of filters.

Although the filterbank tree can be built in a completely general way, some node combinations do not result in well localized frequency bands and changing the subband ordering is not relevant in such cases.

### 3.2. Alternative Boundary Handling

Assuming that the signals are periodic is unnatural in a sense; it may also result in "false" significant wavelet coefficients due to the possible discontinuities. A general

<sup>4</sup>The data structure was not implemented as a Matlab class since Octave in its current version 3.8.2 does not fully support `classdef` yet.

**ALGORITHM 8:** (Near) Best basis selection

**Input:**  $\mathbf{f} \in \mathbb{C}^{L_s}$ , filterbank tree  $\mathcal{T}_{\text{in}}$  built using orthonormal basic filterbanks, cost measure function  $\mathcal{C}$ .

**Output:** Pruned filterbank tree  $\mathcal{T}_{\text{out}}$ , (near) best basis coefficients  $\mathbf{c}$ .

$\mathbf{f}_{\text{norm}} \leftarrow \mathbf{f} / \|\mathbf{f}\|$ ;

$n\text{ToRemove} = \{\}$ ; % Empty set of indices

$\mathbf{c}_{\text{in}}, \text{inIdStack} \leftarrow$  wavelet packet analysis of  $\mathbf{f}_{\text{norm}}$  using Alg. 6 with  $\mathcal{T}_{\text{in}}$ , scaling  $s = 1$ .

**foreach** subband in  $\mathbf{c}_{\text{in}}$  **do**  $C_k \leftarrow \mathcal{C}(\mathbf{c}_{\text{in},k})$ ;

;

$k = -1 + \sum M_n$ ;

**if**  $\mathcal{C}$  is additive **then**

**for**  $n \leftarrow N - 1$  **to** 1 **do**

    % Parent subband index

$id$  **pop** from  $\text{inIdStack}$ ;

    % Combined cost of children subbands

$C = \sum_{m=k-M_n+1}^k C_m$ ;

$k = k - M_n$ ;

**if**  $C_{id} \leq C$  **then**

      add  $n$  to  $n\text{ToRemove}$ ;

**else**

$C_{id} \leftarrow C$ ;

**end**

**end**

**end**

**if**  $\mathcal{C}$  is not additive **then**

  % Empty sets of indices

**for**  $n \leftarrow N - 1$  **to** 1 **do**  $chid_n \leftarrow \{\}$ ;

  ;

**for**  $n \leftarrow N - 1$  **to** 1 **do**

    % Parent subband index

$id \leftarrow$  **pop** from  $\text{inIdStack}$ ;

$chid \leftarrow \{k - M_n + 1, \dots, k\}$ ;

    % Combined cost of children subbands

$C = \mathcal{C}(\mathbf{c}_{chid_n \cup chid})$ ;

$k = k - M_n$ ;

**if**  $C_{id} \leq C$  **then**

      add  $n$  to  $n\text{ToRemove}$ ;

**else**

$chid_{id} \leftarrow chid_n \cup chid$ ;

**end**

**end**

**end**

Delete  $n\text{ToRemove}$  nodes from  $\mathcal{T}_{\text{in}}$  to obtain  $\mathcal{T}_{\text{out}}$ ;

$\mathbf{c} \leftarrow$  filterbank tree analysis of  $\mathbf{f}$  using Alg. 4 with  $\mathcal{T}_{\text{out}}$ ;

approach to handling a general boundary extension is to allow for some additional coefficients which hold information about the boundaries. Such a general case was analyzed in detail in [Rajmic and Průša 2014]. The main idea is to replace the circular convolution with a linear one assuming the signal being extended according to the chosen extension type. This approach however breaks the finite frame abstraction. Other boundary handling tricks can be found in [Taswell and McGill 1994] and elsewhere, but such are not yet included in the toolbox.

*Toolbox conventions.* The computation functions introduced so far accept additional 'zero', 'odd' and 'even' flags for zero, even or odd symmetric extensions.

#### 4. UNDECIMATED WAVELET FRAMES, NORMALIZATION, CANONICAL DUALS

An undecimated overcomplete wavelet frame is obtained by removing the downsampling from a filterbank tree. This makes the representation translation-invariant, but on the other hand highly redundant. Transform of such a type is known under several names: stationary, maximum overlap, redundant and even continuous wavelet transform.

The undecimated transform can be achieved from a filterbank by removing the downsampling operations and upsampling the filters according to the mutirate identity property (see sec. 1.4 and in particular the middle row in Fig. 2). This filterbank

**ALGORITHM 9:** Create a filterbank tree with frequency ordering of subbands

---

**Input:** Filterbank tree  $\mathcal{T}_{\text{in}}$  producing subbands in the natural order.  $N$  nodes of the tree are traversed in a breadth-first order and  $n$ -th node consists of filterbank  $g_m^n, a_m^n$  for  $m \in \{0, 1, \dots, M_n - 1\}$ .

**Output:** Filterbank tree  $\mathcal{T}_{\text{out}}$  producing subbands in the frequency order.

$\mathcal{T}_{\text{out}} \leftarrow \mathcal{T}_{\text{in}}$ ; % The following operations are performed on  $\mathcal{T}_{\text{out}}$ .

Mark all nodes  $n$  as not reordered;

**for**  $n \leftarrow 0$  **to**  $N - 1$  **do**

**if** node  $n$  is reordered and  $M_n$  is odd **then**

$\text{parity} \leftarrow \text{even}$ ;

**else**

$\text{parity} \leftarrow \text{odd}$ ;

**end**

**for**  $m \leftarrow 0$  **to**  $M_n - 1$  **do**

**if** node  $n$  has a next node  $k$  connected at  $m$  and  $m$  is parity **then**

**for**  $p \leftarrow 0$  **to**  $M_k - 1$  **do** % Flip order of filters in node  $k$

$g_p^k \leftarrow g_{M_k - 1 - p}^k$ ;  $a_p^k \leftarrow a_{M_k - 1 - p}^k$ ;

**end**

      Mark node  $n$  as reordered;

**end**

**end**

**end**

---

shape together with a quick evaluation of the convolutions is called the  $\hat{A}$ -trous algorithm [Holschneider et al. 1990].

Since the undecimated transforms use the same filters as their decimated counterparts, scaling in the analysis and/or synthesis step needs to be employed in order to achieve a perfect reconstruction. Three types of scaling are supported: *noscale*, *scale* and *sqr*t causing scaling of each involved filter by 1, by one over the associated subsampling factors or square roots of the previous, respectively. In order to get a perfect reconstruction, *sqr*t must be used in both analysis and synthesis or *noscale* must be used for analysis and *scale* for synthesis, or vice versa.

Algorithms 10 and 11 are valid only for a DWT filterbank tree, but an extension to general filterbanks and to wavelet packets is straightforward and they are also included in LTFAT. The  $\hat{A}$ -trous trick exploits the fact that the upsampled filter in the convolution has only  $L_g$  nonzero samples.

**ALGORITHM 10:**  $\hat{A}$ -trous algorithm analysis

---

**Input:** Input signal  $\mathbf{f} \in \mathbb{C}^L$ , basic wavelet filters  $g_m, a_m$  for  $m \in \{0, 1, \dots, M - 1\}$ , number of levels  $J$ , scaling factors  $s_m \in \{1, 1/\sqrt{a_m}, 1/a_m\}$ .

**Output:** Coefficients  $\mathbf{c}_k$  for  $k \in \{0, 1, \dots, J(M - 1)\}$ .

Scale all filters such that  $g_m \leftarrow s_m g_m$ ;

Continue as in Alg. 2 with lines 5 and 8 changed as follows:

line 5:  $\mathbf{c}_k \leftarrow \mathbf{a} \otimes \left( \uparrow_{a_0^j} g_m \right)$ ;

line 8:  $\mathbf{a} \leftarrow \mathbf{a} \otimes \left( \uparrow_{a_0^j} g_0 \right)$ ;

---

Any undecimated wavelet filterbank tree or packet can be equally represented as a simple undecimated filterbank which, in turn, can be regarded as a painless, frequency domain version of the nonstationary Gabor transform [Balazs et al. 2011]. Con-

**ALGORITHM 11:** À-trous algorithm synthesis

**Input:** Coefficients  $c_k$  for  $k \in \{0, 1, \dots, J(M-1)\}$ , basic wavelet filters  $g_m, a_m$  for  $k \in \{0, 1, \dots, M-1\}$ , number of levels  $J$ , scaling factors  $s_m \in \{1/a_m, 1/\sqrt{a_m}, 1\}$ .

**Output:** Reconstructed signal  $\mathbf{f} \in \mathbb{C}^L$ .

Scale all filters such that  $g_m \leftarrow s_m g_m$ ;

Continue as in Alg. 3 with lines 4 and 6 changed as follows;

line 4:  $\mathbf{a} \leftarrow \mathbf{a} \otimes \left( \uparrow_{a_0^j} g_0 \right)$ ;

line 6:  $\mathbf{a} \leftarrow \mathbf{a} + \mathbf{c}_k \otimes \left( \uparrow_{a_0^j} g_m \right)$ ;

sequently, the canonical dual frame can be obtained easily by exploiting the fact that the frame operator is diagonal in the frequency domain.

The downside of this approach is that the canonical dual frame will not keep the iterated filterbank structure, but it becomes a uniform filterbank which admits a potentially less effective algorithm. In general, Alg. 12 produces the canonical dual frame valid for a single length  $L$  since the support of the dual filters might not be compact any more. A reasonable FIR approximation can usually be obtained through truncation.

**ALGORITHM 12:** Find the canonical dual frame of an undecimated wavelet frame

**Input:** General undecimated filterbank tree  $\mathcal{T}_{\text{in}}$  with  $K$  outputs, system length  $L$ , frequency ordering boolean flag (*freqOrdering*).

**Output:** Dual frame undecimated filterbank  $\mathbf{h}_k, a = 1$  for  $k \in \{0, 1, \dots, K-1\}$  for a system length  $L$ .

**if** *freqOrdering* **then**  $\mathcal{T}_{\text{in}} \leftarrow$  Alg. 9 with  $\mathcal{T}_{\text{in}}$ ;

;

**foreach** *k-th output of*  $\mathcal{T}_{\text{in}}$  **do**

$g_k \leftarrow$  identical filter using Alg. 1;

$\widehat{\mathbf{g}}_k \leftarrow \mathcal{F}^*(g_k(L))$ ;

**end**

**for**  $k \leftarrow 0$  **to**  $K-1$  **do**

$\mathbf{h}_k \leftarrow \mathcal{F}\left(\widehat{\mathbf{g}}_k / \sum_{m=0}^{M-1} |\widehat{\mathbf{g}}_k|^2\right)$ ; % Squaring and division are performed element-wise.

**end**

*Toolbox conventions:* The functions performing the undecimated versions of transforms from the previous sections are: `ufwt`, `iufwt`, `uwfbt`, `iuwfbt`, `uwpfbt` and `iuwpfbt`. The canonical dual frame can be computed using a combination of functions `wfbt2filterbank` or `wpfbt2filterbank` and `filterbankdual`.

**5. DUAL-TREE COMPLEX WAVELET TRANSFORM**

The dual-tree complex wavelet transform (DT-CWT) is a two-times redundant representation based on a pair of DWT filterbank trees behaving as a Hilbert transform pair. In the original design [Kingsbury 2000; 2001; 2003; Selesnick et al. 2005], all the filters are FIR and the Hilbert transform relationship is only approximate since the exact relationship would enforce infinitely supported filters in one of the trees or a pre-processing with an infinitely supported Hilbert transformer. If the coefficients are properly combined, the overall transform is almost analytic, meaning that the effective frame vectors are complex and their Fourier transform is supported largely in the positive or negative frequency domains.

DT-CWT has the advantage over the ordinary DWT that the subbands are almost aliasing-free, which makes the representation near translation-invariant and more robust to undesired effects during the coefficient processing. Moreover, the real input signals require only half of the coefficients to be computed since the other half is just a complex conjugate, much like the DFT of a real signal.

The original design procedure suggests to use one sample delay difference in the very first level of trees which results in a better analytic behaviour of the first levels of the transform. In fact, any basic wavelet filterbank can be used in the first level of both trees, as long as it has the same number of filters and equal downsampling rates as the rest of the tree and provided that the filters are shifted by one sample in one of the trees. In effect, such trees have to be treated as general filterbank trees as described in sec. 3.

A natural extension was proposed in [Selesnick 2004], where the dual-density and dual-tree wavelet transforms were combined. In such a case, the tree nodes are three-channel filterbanks and the algorithms still apply.

The coefficients of a dual tree complex wavelet transform can be obtained from a signal by Alg. 13 and the synthesis operation is performed by Alg. 14.

As in the case of general filterbank trees, the overall dual-tree filterbank can be equally represented by a non-iterated filterbank obtained by combining multirate identity representations of both trees.

---

**ALGORITHM 13: Dual-Tree complex wavelet analysis**


---

**Input:** Input signal  $\mathbf{f} \in \mathbb{C}^{L_s}$ , Hilbert transform pair of basic filterbanks  $g_{a,m}$  and  $g_{b,m}$ ,  $a_m$ , first level filterbank  $g_m$ ,  $m \in \{0, 1, \dots, M-1\}$ , number of levels  $J$ .

**Output:** Complex coefficients  $\mathbf{c}_k$  for  $k \in \{0, 1, \dots, 2K-1\}$  where  $K = J(M-1) + 1$ .

- 1 Form a DWT filterbank tree  $\mathcal{T}_a$  using  $g_{a,m}$  and  $g_m$  in the first level;
  - 2 Form a DWT filterbank tree  $\mathcal{T}_b$  using  $g_{b,m}$  and  $g_m$  delayed by a 1 sample in the first level;
  - 3  $\mathbf{c}_{a,k} \leftarrow$  Alg. 4 with  $\mathcal{T}_a$ ;  $\mathbf{c}_{b,k} \leftarrow$  Alg. 4 with  $\mathcal{T}_b$ ;
  - 4 **for**  $k \leftarrow 0$  **to**  $K-1$  **do**
  - 5      $\mathbf{c}_k \leftarrow \frac{1}{2}(\mathbf{c}_{a,k} + i\mathbf{c}_{b,k})$ ;  $\mathbf{c}_{2K-1-k} \leftarrow \frac{1}{2}(\mathbf{c}_{a,k} - i\mathbf{c}_{b,k})$ ;
  - % Subbands  $\mathbf{c}_k$  and  $\mathbf{c}_{2K-1-k}$  are complex conjugates if  $\mathbf{f} \in \mathbb{R}^L$
  - 6 **end**
- 

---

**ALGORITHM 14: Dual-Tree complex wavelet synthesis**


---

**Input:** Coefficients  $\mathbf{c}_k$  for  $k \in \{0, 1, \dots, 2K-1\}$  where  $K = J(M-1) + 1$ , Hilbert transform pair of basic filterbanks  $g_{a,m}$  and  $g_{b,m}$ ,  $a_m$ , first level filterbank  $g_m$ ,  $m \in \{0, 1, \dots, M-1\}$ , number of levels  $J$ .

**Output:** Reconstructed signal  $\tilde{\mathbf{f}} \in \mathbb{C}^{L_s}$ .

- 1 Form a DWT filterbank tree  $\mathcal{T}_a$  using  $g_{a,m}$  and  $g_m$  in the first level;
  - 2 Form a DWT filterbank tree  $\mathcal{T}_b$  using  $g_{b,m}$  and  $g_m$  delayed by a 1 sample in the first level;
  - 3 **for**  $k \leftarrow 0$  **to**  $K-1$  **do**
  - 4      $\mathbf{c}_{a,k} \leftarrow \frac{1}{2}(\mathbf{c}_k + \mathbf{c}_{2K-1-k})$ ;  $\mathbf{c}_{b,k} \leftarrow i\frac{1}{2}(\mathbf{c}_{2K-1-k} - \mathbf{c}_k)$ ;
  - %  $\mathbf{c}_{a,k} \leftarrow \text{real}(\mathbf{c}_k)$ ;  $\mathbf{c}_{b,k} \leftarrow \text{imag}(\mathbf{c}_k)$  if  $\mathbf{f} \in \mathbb{R}^L$
  - 5 **end**
  - 6  $\tilde{\mathbf{f}}_a \leftarrow$  Alg. 5 with  $\mathcal{T}_a$  and  $\mathbf{c}_{a,k}$ ;  $\tilde{\mathbf{f}}_b \leftarrow$  Alg. 5 with  $\mathcal{T}_b$  and  $\mathbf{c}_{b,k}$ ;  $\tilde{\mathbf{f}} \leftarrow \tilde{\mathbf{f}}_a + \tilde{\mathbf{f}}_b$ ;
-



*Toolbox Conventions.* Functions implementing Alg. 13 are called `dtwfb` and `dtwfbreal` for real-only signals. Similarly, there are functions `idtwfb` and `idtwfbreal` implementing Alg. 14.

Hilbert transform filterbank pairs suitable for the second and higher levels of trees are stored in functions with `wfiltdt_` prefix: odd and even length biorthogonal filterbanks `oddeven` [Kingsbury 2001], quarter-shift orthogonal filterbanks `qshift` [Kingsbury 2000; 2003], optimized symmetric `optsym` [Dumitrescu et al. 2008] and double-density dual-tree filterbanks `dden` [Selesnick 2004].

## 6. GENERAL DUAL-TREE COMPLEX FILTERBANK

A generalization of a dual-tree complex wavelet transform analogue to the generalization of DWT described in sec. 3 was proposed in [Bayram and Selesnick 2008]. It is not straightforward as it requires the high-pass filters in the basic filterbanks to meet the half-sample shift condition as well and it further requires another set of *extension* basic wavelet filters to be used in specific nodes in the tree in order to preserve the analytic behaviour. The extension filterbank is shared by both the trees and any basic wavelet filterbank can be used provided it has equal number of filters and downsampling factors as the other filterbanks. Algorithms for analysis and synthesis are Alg. 15 and Alg. 16 respectively.

---

### ALGORITHM 15: General dual-tree complex wavelet filterbank analysis

---

**Input:** Input signal  $f \in \mathbb{C}^{L_s}$ , Hilbert transform pair of basic filterbanks  $g_{a,m}$  and  $g_{b,m}$ ,  $a_m$ , first level filterbank  $g_{root,m}$  and extension filterbank  $g_{ext,m}$ ,  $m \in \{0, 1, \dots, M-1\}$ , shape of a tree  $\mathcal{T}_{shape}$ .

**Output:** Complex coefficients  $c_k$  for  $k \in \{0, 1, \dots, 2K-1\}$  where  $K$  is the total number of unconnected outputs of one of the trees.

Form a filterbank tree  $\mathcal{T}_a$  of a shape  $\mathcal{T}_{shape}$  using  $g_{root,m}$  in the root node and using  $g_{a,m}$  in its direct descendants and in their further lowpass filter connections and using  $g_{ext,m}$  elsewhere; Create  $\mathcal{T}_b$  in the same manner using  $\mathcal{T}_b$  and one sample delayed version of  $g_{root,m}$ ; Continue as in Alg. 13 starting with line 3;

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### ALGORITHM 16: General dual-tree complex wavelet filterbank synthesis

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**Input:** Coefficients  $c_k$  for  $k \in \{0, 1, \dots, 2K-1\}$  where  $K$  is the total number of unconnected outputs of one of the trees, Hilbert transform pair of basic filterbanks  $g_{a,m}$  and  $g_{b,m}$ ,  $a_m$ , first level filterbank  $g_{root,m}$  and extension filterbank  $g_{ext,m}$ ,  $m \in \{0, 1, \dots, M-1\}$ , shape of a tree  $\mathcal{T}_{shape}$ .

**Output:** Reconstructed signal  $\tilde{f} \in \mathbb{C}^{L_s}$ .

Form a filterbank tree  $\mathcal{T}_a$  of a shape  $\mathcal{T}_{shape}$  using  $g_{root,m}$  in the root node and using  $g_{a,m}$  in its direct descendants and in their further lowpass filter connections and using  $g_{ext,m}$  elsewhere; Create  $\mathcal{T}_b$  in the same manner using  $\mathcal{T}_b$  and one sample delayed version of  $g_{root,m}$ ; Continue as in Alg. 14 starting with line 3;

---

*Toolbox conventions:* General dual-tree complex filterbanks are supported directly by functions `dtwfb*` and `idtwfb*`. The available predefined shapes are as in sec. 3 and the tree structure can be created using `dtwfbinit` and pruned by `wfbtremove`. Adding nodes using `wfbtput` is not yet supported.

## 7. FRAME BOUNDS OF GENERAL FILTERBANK TREES

This section describes a computationally tractable algorithm for calculating optimal frame bounds  $A, B$  of a wavelet frame (any type described in this paper) for large  $L$  while avoiding explicit computation of singular values of  $F$  (or equivalently eigenvalues of  $FF^*$ ).

The ratio  $\sqrt{B/A}$  gives the condition number of a matrix  $F$  with frame vectors as columns, which determines the numeric stability of the system, and  $B$  alone corresponds to the spectral (operator) norm of such a matrix, as  $\|F\| = \sqrt{B}$ . Performance of many iterative algorithms depends on the knowledge of frame bounds and/or they are restricted to be working with matrices  $\|F\| \leq 1$  only (e.g. the frame algorithm [Gröchenig 1993], iterative soft thresholding [Daubechies et al. 2004] and proximal convex optimization algorithms in general [Combettes and Pesquet 2011]).

The literature on the frame bounds of iterated filterbanks is sparse, and only the DWT case in  $\ell^2(\mathbb{Z})$  was addressed as the analytical treatment of general filterbank trees would become cumbersome. In [Stanhill and Zeevi 1996], the authors established some connection between the bounds of the basic wavelet filterbank and the  $J$ -level DWT. Namely the following holds for a basic filterbank with bounds  $A_1, B_1 \geq 1$

$$A_J \geq A_1 \text{ and } B_J \leq B_1^J, \quad (5)$$

where the subscripts denote number of levels in the DWT filterbank. For tight basic filterbanks with bounds equal to  $A_1 = B_1 > 1$  the following was proven

$$A_J = A_1 \text{ and } B_J \leq A_1^J, \quad (6)$$

which shows that the tightness might be lost during iterations. The exception is the Parseval tightness  $A_1 = 1$ , which is preserved under iterations  $A_J = B_J = 1$ . Results (5),(6) also hold for the undecimated DWT. It has been observed in [Fowler 2005] that the bounds of the  $J$ -level undecimated DWT with (originally orthonormal) two-channel basic filterbank with  $A_1 = B_1 = 2$  (noscale flag from sec. 4) are  $A_J = 2, B_J = 2^J$  exactly.

We observed the same behaviour for bounds for frames in  $\mathbb{C}^L$ , calculated using algorithms described in this section.

Although the precise relationship between the frame bounds in  $\mathbb{C}^L$  and  $\ell^2(\mathbb{Z})$  has not yet been addressed in the literature, in the practical applications conducted in  $\mathbb{C}^L$ , we are usually interested in bounds of the frame at hand.

The numerical evaluation of frame bounds of a general filterbank tree is not straightforward and it requires two pre-processing steps:

- (1) Using properties from sec. 1.4, transform a filterbank tree of interest to an identical, possibly non-uniform filterbank.
- (2) Transform the resulting filterbank to an uniform one using the identity by [Akkarakaran and Vaidyanathan 2003] (summarized in Alg. 17).

Having an equivalent uniform filterbank frame, the optimal frame bounds in  $\mathbb{C}^L$  can be obtained such that the lower frame bound is the the minimum of lower frame bounds of a length- $\frac{L}{a}$  series of smaller frames (see below) in  $\mathbb{C}^a$  and, likewise, the upper frame bound is the maximum of the upper frame bounds of such frames [Fickus et al. 2013]. Although not specifically mentioned in their paper, such a behaviour can be explained through the analysis of the frame operator, which exhibits a regular behaviour in both the time and the frequency domains. In the time domain, the frame operator is a block-circulant matrix and in the frequency domain, it can become a block-diagonal matrix consisting of  $\frac{L}{a}$  blocks of size  $a \times a$  after suitable row and column permutations. Eigenvalues of a block-diagonal matrix are just the concatenation of eigenvalues of separate

block-matrices and therefore the frame bounds are easily obtained. For sufficiently small  $a$ , the blocks can be created explicitly and the eigenvalues of blocks evaluated directly as shown in Alg. 18.

The block-diagonal form of the frame operator can be also exploited for effective computation of canonical dual frames of uniform filterbanks and even nonuniform ones with  $\text{lcm}(a_0, a_1, \dots, a_{M-1})$  being reasonably low. We omit the description of the algorithm as it is not the main objective of the paper and we refer the interested reader to function `filterbankdual` and `nonu2filterbank`.

---

**ALGORITHM 17: Non-uniform to uniform filterbank transformation**


---

**Input:** Nonuniform filterbank  $g_m, a_m$  for  $m \in \{0, 1, \dots, M-1\}$ .

**Output:** Uniform filterbank  $h_k, a$  for  $k \in \{0, 1, \dots, K-1\}$ , where

$$K = \sum_{m=0}^{M-1} \text{lcm}(a_0, a_1, \dots, a_{M-1}) / a_m.$$

$a = \text{lcm}(a_0, a_1, \dots, a_{M-1});$

$k = 0;$

**for**  $m \leftarrow 0$  **to**  $M-1$  **do**

**for**  $p \leftarrow 0$  **to**  $a/a_m - 1$  **do**

$h_k \leftarrow g_m$  with  $d_{h_k} = d_{g_m} - a_m p$ ; % Adjust filter ‘offset’

$k \leftarrow k + 1;$

**end**

**end**

---



---

**ALGORITHM 18: Frame bounds of an uniform filterbank in  $\mathbb{C}^L$** 


---

**Input:** Uniform filterbank  $g_m \in \mathbb{C}^L$  and  $a$  for  $m \in \{0, 1, \dots, M-1\}$ , system length  $L = Na$  for some  $N \in \mathbb{N}$ .

**Output:** Optimal frame bounds  $A, B$ .

$A \leftarrow \infty, B \leftarrow 0;$

$\mathbf{G} \leftarrow \text{stack } \widehat{\mathbf{g}}_m \text{ as columns};$  %  $\mathbf{G} \in \mathbb{C}^{L \times M}$

**for**  $w \leftarrow 0$  **to**  $N-1$  **do**

$\mathbf{H} \leftarrow \text{stack rows } (w + kN) \text{ of } \mathbf{G} \text{ for } k \in \{0, 1, \dots, a-1\};$  %  $\mathbf{H} \in \mathbb{C}^{a \times M}$

$\mathbf{e} \leftarrow \text{eig}(\mathbf{H}\mathbf{H}^*);$  % Compute eigenvalues of a  $a \times a$  matrix

$A \leftarrow \min(A, \min(\mathbf{e})); B \leftarrow \max(B, \max(\mathbf{e}));$

**end**

$A \leftarrow A/a; B \leftarrow B/a;$

---

Alg. 18 can be simplified for undecimated transforms ( $a = 1$ ) while the frame operator in the Fourier domain is diagonal (Alg. 19) (and even constant in the tight frame case).

---

**ALGORITHM 19: Frame bounds of an undecimated filterbank in  $\mathbb{C}^L$** 


---

**Input:** Undecimated filterbank  $g_m \in \mathbb{C}^L$  for  $m \in \{0, 1, \dots, M-1\}$ , system length  $L \in \mathbb{N}^+$ .

**Output:** Optimal frame bounds  $A, B$ .

$\mathbf{h} \leftarrow \sum_{m=0}^{M-1} |\widehat{\mathbf{g}}_m|^2;$  % Square and the abs. operation are performed element-wise.

$A \leftarrow \min(\mathbf{h}); B \leftarrow \max(\mathbf{h});$

---

*Toolbox conventions.* Toolbox functions for calculating frame bounds are called `wfbbounds` and `wpfbbounds` for both the decimated case and for the undecimated transforms. Note there are no such functions for `fwf` and `ufwf` as they are both special cases of `wf` and `wf` respectively. The function calculating the frame bounds of the dual-tree wavelet transform from sec. 5 is called `dtwfbbounds`. All functions require defining compatible  $L$ , equal or longer than the longest identical filter according to Alg. 1 and flags specific to the respective transforms.

## 8. CONCLUSION AND OUTLOOK

The hope of the authors is that the toolbox becomes a base for new numerical experiments and scientific developments.

In the future, apart from the topics already mentioned in the introduction, the toolbox might be enriched by inclusion of the following:

Support of a systematic construction of tight wavelet frames and bi-framelets [Ron and Shen 1997; Benedetto and Li 1998; Daubechies et al. 2003; Ehler 2007] through unitary, oblique and mixed extension principles, which is an active area of research even nowadays. As the fast algorithms mostly coincide with the already implemented fast wavelet transform with more than two wavelet filters, it is expected that the results will fit into the already existing framework seamlessly.

Inclusion of more general wavelet frames (not based on MRA) via frequency warping [Holighaus et al. 2015; Christensen and Goh 2014]. These recent results might be the answer to the ongoing quest for the truly discrete wavelet frames with a denser logarithmic frequency resolution.

## APPENDIX

This section contains Alg. 20 and Fig. 5 which were moved here to avoid cluttering Section 3. Alg. 20 is used as a subprocess in Alg. 4 and Alg. 5.

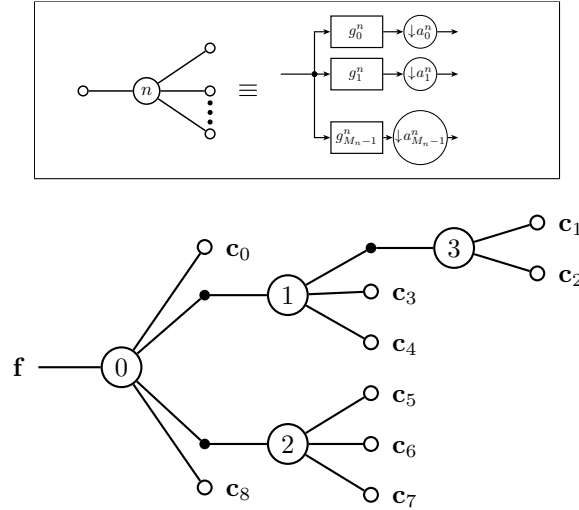


Fig. 5. An example of a schematic diagram of a filterbank tree as used in Alg. 4, Alg. 5 and Alg. 20. The tree consists of  $N = 4$  nodes, the number of filters  $M_n$  in node  $n$  is  $M_0 = 4$ ,  $M_1 = 3$ ,  $M_2 = 3$ ,  $M_3 = 2$ , where  $n$  (the number in the nodes) is the index in the breadth-first order. There are  $K = 9$  unconnected outputs of the tree shown as empty circles. The output of Alg. 20 for this particular tree shape is  $kIdMap = \{0, 8, 3, 4, 5, 6, 7, 1, 2\}$ .

---

**ALGORITHM 20:** Compute filterbank tree output index map

---

**Input:** Filterbank tree  $\mathcal{T}$  with  $N$  nodes in the breadth-first order where the  $n$ -th node consists of  $M_n$  filters and  $M_{\text{unc},n}$  of them are unconnected.

**Output:** Array of indices  $kIdMap$  of length  $K = \sum_{n=0}^{N-1} M_{\text{unc},n}$  (the total number of unconnected outputs of all nodes).

```

 $k \leftarrow 0$  ;
Prepare array  $kIdMap$  of length  $K$  ;
 $nodeFiltTuple \leftarrow \{\}$  ; % an empty LIFO stack of index tuple  $(n, m_{\text{start}}, m_{\text{unc}})$ 
push  $(0, 0, 0)$  to  $nodeFiltTuple$  ;
while  $nodeFiltTuple$  is not empty do
   $(n, m_{\text{start}}, m_{\text{unc}}) \leftarrow \text{pop from } nodeFiltTuple$  ;
  for  $m \leftarrow m_{\text{start}}$  to  $M_n - 1$  do
    if node  $n$  has a next node  $n_{\text{next}}$  connected at index  $m$  then
      if  $m < M_n - 1$  then
        push  $(n, m + 1, m_{\text{unc}})$  to  $nodeFiltTuple$  ;
      end
      push  $(n_{\text{next}}, 0, 0)$  to  $nodeFiltTuple$  ;
      break ;
    else
       $kIdMap(m_{\text{unc}} + \sum_{l=0}^{n-1} M_{\text{unc},l}) = k$  ;
       $k \leftarrow k + 1$  ;  $m_{\text{unc}} \leftarrow m_{\text{unc}} + 1$  ;
    end
  end
end

```

---

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